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Privacy Preserving Aggregation of Distributed Mobility Data Streams



Thomas Liebig

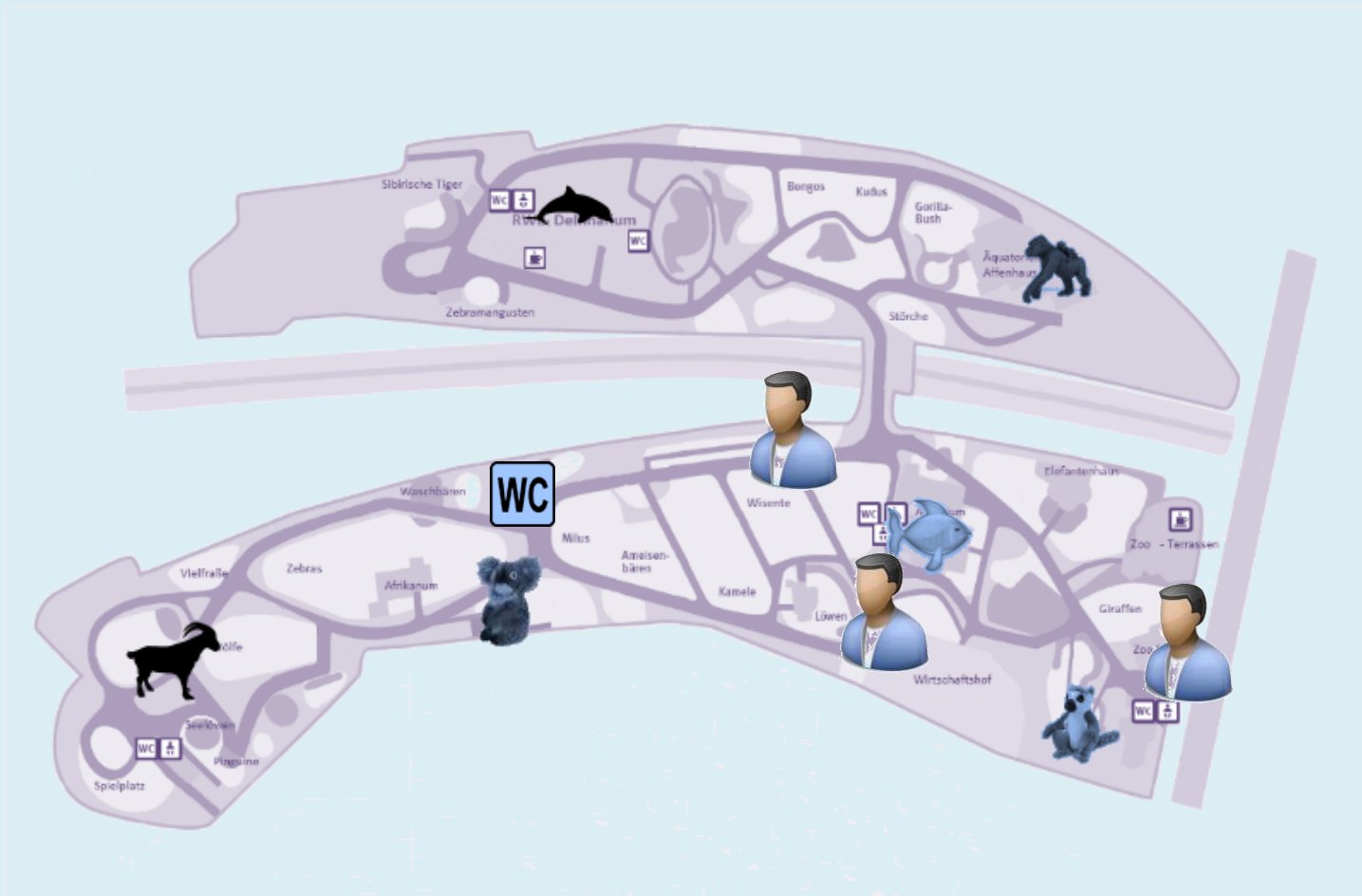
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Mobility Data Aggregation Problem



□ How many people have been at which location?



Re-Identification problems



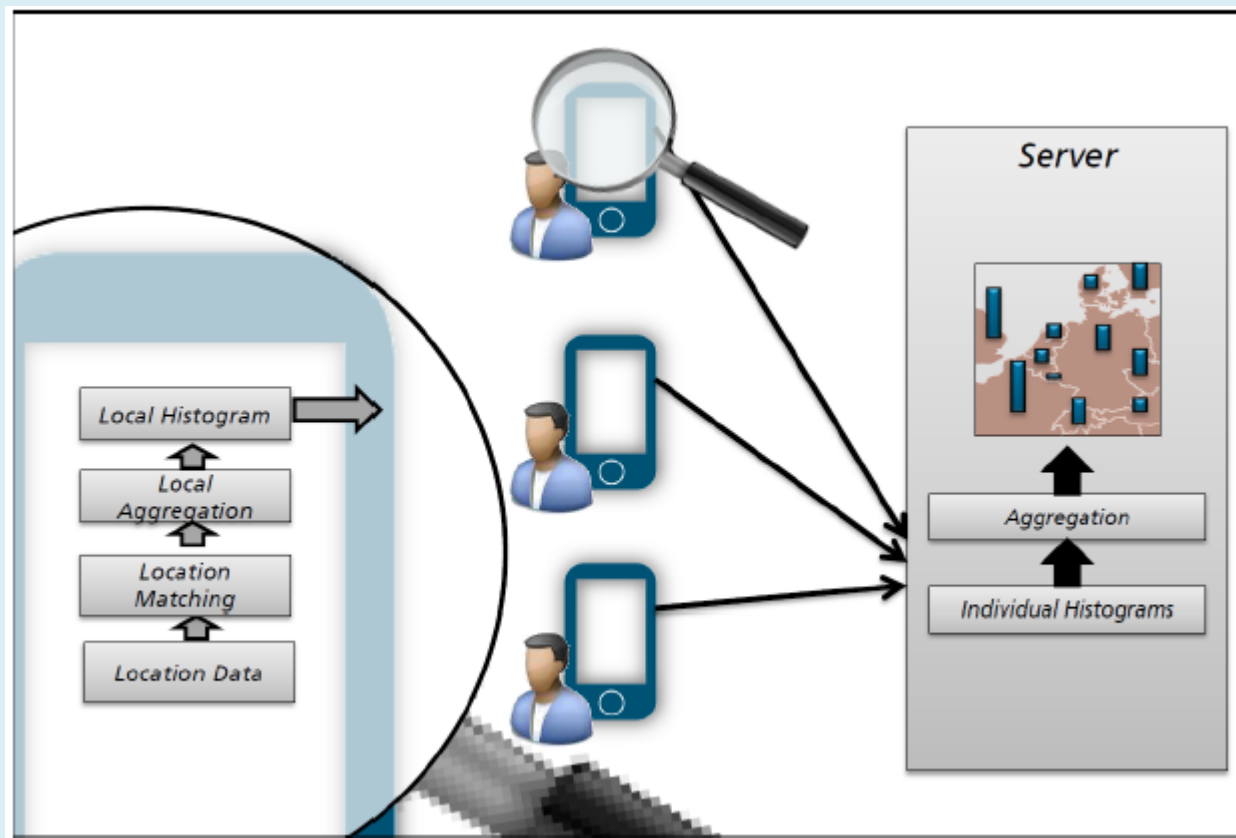
- removing personal identifier insufficient
 - sending position information in bulks (trajectories or histograms) insufficient
 - Easy re-identification in mobility data
 - Intrusted server
- streaming solution required which prevents server from accessing individual mobility traces

Problem described in [Andrienko et. al. 12] (with many examples and references to laws)

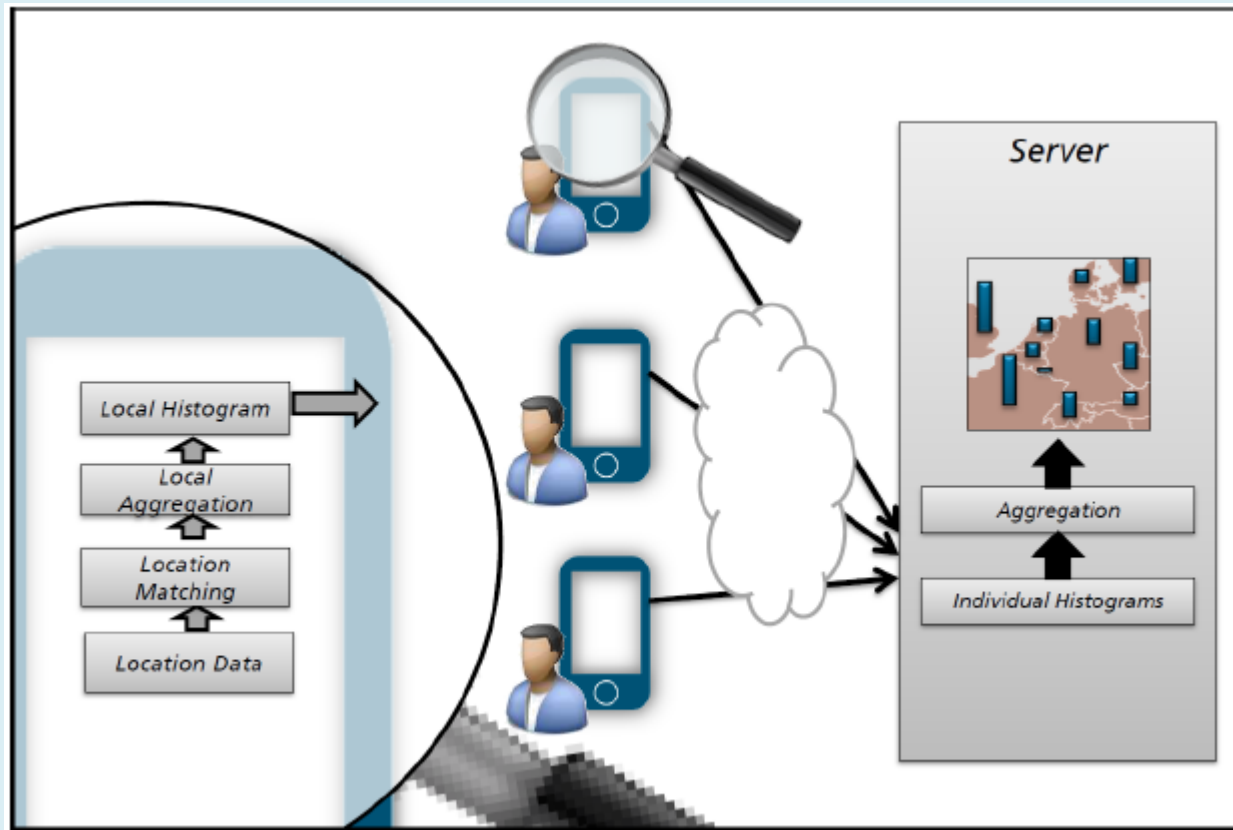
Problem



□ secure this aggregation:



Related approach [Kopp et. al. 12]



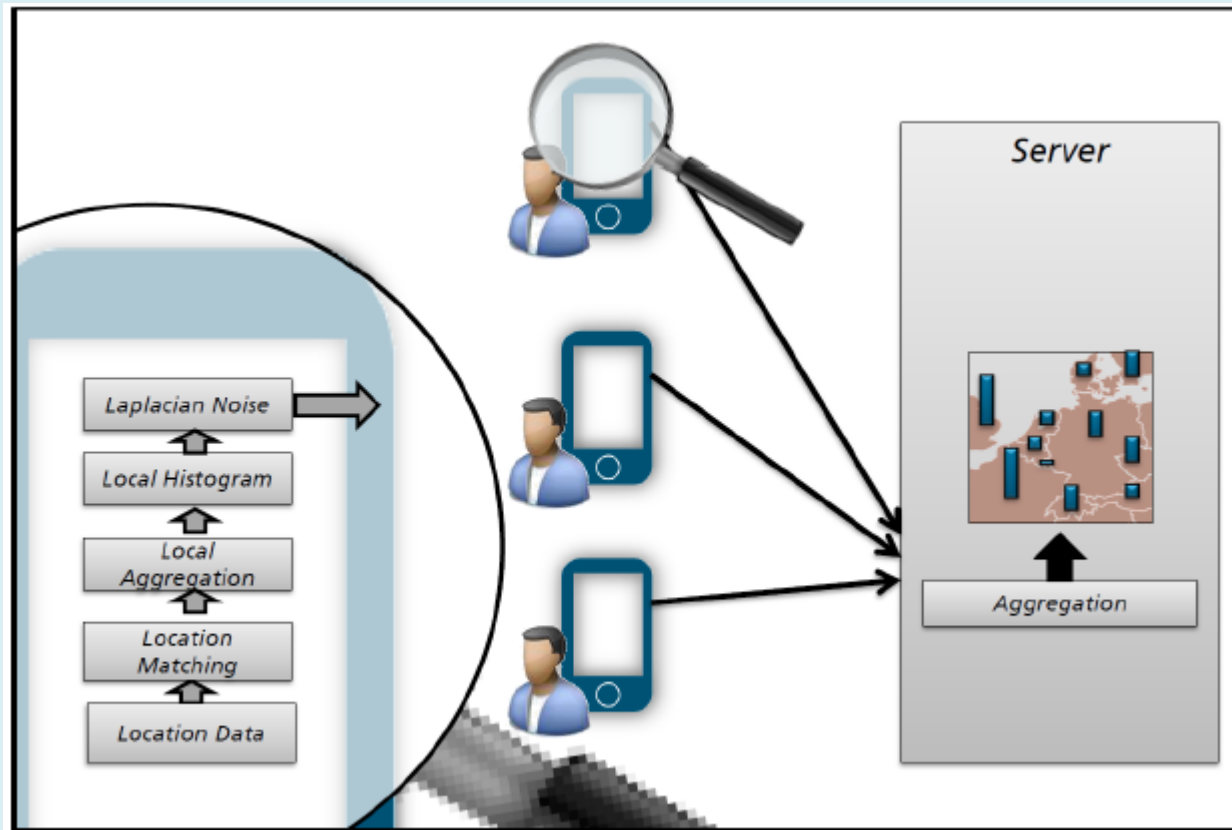
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Problem described in [Andrienko et. al. 12] (with many examples and references to laws)

Related approach [Monreale et. al. 13]



Re-Identification problems



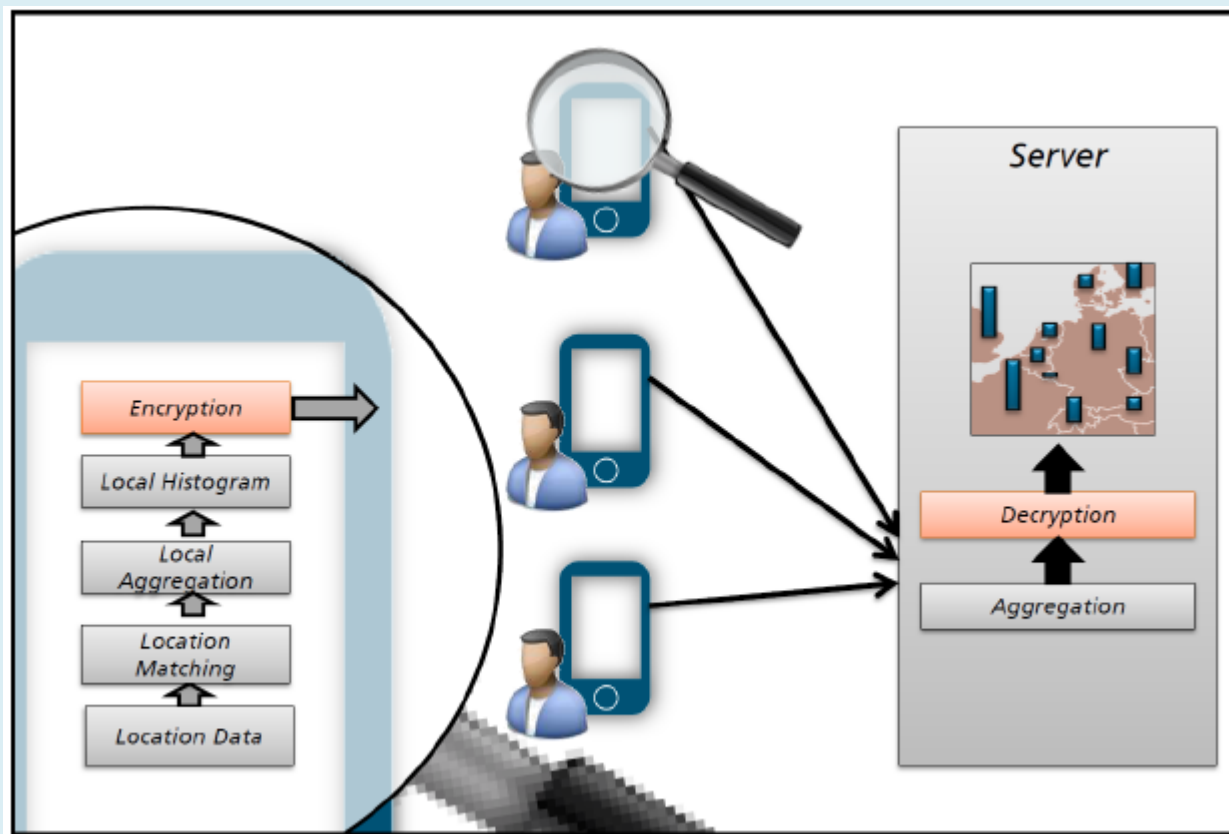
- removing personal identifier insufficient
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Proposed Solution (3 steps)



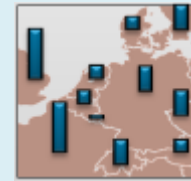
- utilize homomorphic encryption



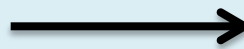
Proposed Solution I (RSA)



- based on two cryptographic keys (like in ssh)



$M := \text{Encryption}(H, \text{public key})$



$\text{Decryption}(M, \text{private key})$

$=$
 H

Proposed Solution II (Damgaard Algorithm)



□ sharing the keys (homomorphic cryptography)



$M1 := \text{Encryption}(H1, \text{public key1})$

$M2 := \text{Encryption}(H2, \text{public key2})$

$M3 := \text{Encryption}(H3, \text{public key3})$



$\text{Decryption}(M1 * M2 * M3, \text{private key})$
=
 $H1 + H2 + H3$

Proposed Solution II (Hash Chain)



□ prevent timing attacks

Proposed Solution



- M1:=Encryption (H1*Ti, public key1)
- M2:=Encryption (H2*Ti, public key2)
- M3:=Encryption (H3*Ti, public key3)



$$\text{Decryption (M1*M2*M3, private key)} \\ = \\ (H1+H2+H3)*Ti$$

Summary



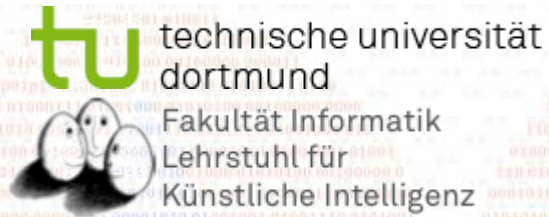
- privacy-by-design concept for mobility stream aggregation
- full utility
- no randomization, no approximation

Next steps:

- further algorithms (clustering, classification, subgroup discovery) without revealing the data

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